

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

DATA STRUCTURE USING C LABORATORY		SEMESTER	I
Course Code	BCAL208	CIE Marks	50
Teaching Hours/Week (L: P: SDA)	1:2:0	SEE Marks	50
Total Hours of Pedagogy	14 Sessions	Total Marks	100
Credits	02	Exam Hours	03
Type of Course SEE	Practical		
COURSE OBJECTIVES: <ul style="list-style-type: none">Equip students with the skills to implement and manipulate fundamental data structures like arrays, stacks, queues, and linked lists.Enable students to solve complex problems using recursion, matrix operations, and tree traversal techniques.			

PART-A

1. Write a program to implement and demonstrate the following operations on an array: Insertion, Deletion, Traversing, and Searching.

Algorithm.**Array - Insertion Operation:**

1. Start

2. Create an Array of a desired datatype and size.

3. Initialize a variable 'i' as 0.

4. Enter the element at ith index of the array.

5. Increment i by 1.

6. Repeat Steps 4 & 5 until the end of the array.

7. Stop

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Array - Deletion Operation:

1. Start
2. Set $J = K$
3. Repeat steps 4 and 5 while $J < N$
4. Set $LA[J] = LA[J + 1]$
5. Set $J = J + 1$
6. Set $N = N - 1$
7. Stop

Array - Search Operation:

1. Start
2. Set $J = 0$
3. Repeat steps 4 and 5 while $J < N$
4. IF $LA[J]$ is equal ITEM THEN GOTO STEP 6
5. Set $J = J + 1$
6. PRINT $J, ITEM$
7. Stop

Array - Traversal Operation:

- 1 Start
2. Initialize an Array of certain size and datatype.
3. Initialize another variable i with 0.
4. Print the i th value in the array and increment i .
5. Repeat Step 4 until the end of the array is reached.
6. End

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Program:

```
#include <stdio.h>
```

```
#define SIZE 5
```

```
void display(int arr[], int n) {  
    printf("Array: ");  
    for(int i=0; i<n; i++) printf("%d ", arr[i]);  
    printf("\n");  
}
```

```
int search(int arr[], int n, int key) {  
    for(int i=0; i<n; i++)  
        if(arr[i] == key) return i;  
    return -1;  
}
```

```
void insert(int arr[], int *n, int pos, int val) {  
    if(*n >= SIZE) {  
        printf("Overflow!\n");  
        return;  
    }  
    for(int i=*n; i>pos; i--) arr[i] = arr[i-1];  
    arr[pos] = val;  
    (*n)++;  
}
```

```
void delete(int arr[], int *n, int pos) {  
    if(*n <= 0) {
```

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```
        printf("Underflow!\n");
        return;
    }
    for(int i=pos; i<*n-1; i++) arr[i] = arr[i+1];
    (*n)--;
}

int main() {
    int arr[SIZE] = {1, 2, 3}, n = 3;
    display(arr, n);
    insert(arr, &n, 1, 5);
    printf("After insertion: ");
    display(arr, n);

    delete(arr, &n, 2);
    printf("After deletion: ");
    display(arr, n);

    int key = 5;
    int pos = search(arr, n, key);
    if(pos != -1) printf("%d found at index %d\n", key, pos);
    else printf("%d not found\n", key);
    return 0;
}
```

Output:

Array: 1 2 3

After insertion: Array: 1 5 2 3

After deletion: Array: 1 5 3

5 found at index 1

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

2. Write a recursive program to solve Towers of Hanoi problem for n disks

Algorithm.

To write an algorithm for Tower of Hanoi, first we need to learn how to solve this problem with lesser amount of disks, say \rightarrow 1 or 2. We mark three towers with name, **source**, **destination** and **aux** (only to help moving the disks). If we have only one disk, then it can easily be moved from source to destination peg.

If we have 2 disks –

- First, we move the smaller (top) disk to aux peg.
- Then, we move the larger (bottom) disk to destination peg.
- And finally, we move the smaller disk from aux to destination peg.

So now, we are in a position to design an algorithm for Tower of Hanoi with more than two disks. We divide the stack of disks in two parts. The largest disk (n^{th} disk) is in one part and all other ($n-1$) disks are in the second part.

Our ultimate aim is to move disk **n** from source to destination and then put all other ($n-1$) disks onto it. We can imagine to apply the same in a recursive way for all given set of disks.

The steps to follow are –

```
Step 1 - Move n-1 disks from source to aux
Step 2 - Move nth disk from source to dest
Step 3 - Move n-1 disks from aux to dest
```

A recursive algorithm for Tower of Hanoi can be driven as follows –

```
START
```

```
Procedure Hanoi(disk, source, dest, aux)
```

```
    IF disk == 1, THEN
```

```
        move disk from source to dest
```

```
    ELSE
```

```
        Hanoi(disk - 1, source, aux, dest) // Step 1
```

```
        move disk from source to dest // Step 2
```

```
        Hanoi(disk - 1, aux, dest, source) // Step 3
```

```
    END IF
```

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

END Procedure

STOP

Program.

```
#include <stdio.h>

void towers(int n, char from, char to, char aux) {
    if(n == 1) {
        printf("Move disk 1 from %c to %c\n", from, to);
        return;
    }
    towers(n-1, from, aux, to);
    printf("Move disk %d from %c to %c\n", n, from, to);
    towers(n-1, aux, to, from);
}

int main() {
    int n = 3;
    printf("Tower of Hanoi solution for %d disks:\n", n);
    towers(n, 'A', 'C', 'B');
    return 0;
}
```

Output:

Tower of Hanoi solution for 3 disks:

Move disk 1 from A to C

Move disk 2 from A to B

Move disk 1 from C to B

Move disk 3 from A to C

Move disk 1 from B to A

Move disk 2 from B to C

Move disk 1 from A to C

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Explanation:

The recursive solution moves n disks from source peg to destination peg using an auxiliary peg:

1. Move $n-1$ disks from source to auxiliary peg
2. Move n th disk from source to destination
3. Move $n-1$ disks from auxiliary to destination peg

3. Write a recursive program to calculate Greatest Common Divisor (GCD) of two numbers.

Algorithm.

Refer an algorithm given below to find the greatest common divisor (GCD) for the given two numbers by using the recursive function.

Step 1 – Define the recursive function.

Step 2 – Read the two integers a and b .

Step 3 – Call recursive function.

- a. if $i > j$
- b. then return the function with parameters i, j
- c. if $i == 0$
- d. then return j
- e. else return the function with parameters $i, j \% i$.

Program:

```
#include <stdio.h>

int gcd(int a, int b) {
    if(b == 0) return a;
    return gcd(b, a % b);
}

int main() {
    int a = 48, b = 18;
    printf("GCD of %d and %d is %d\n", a, b, gcd(a, b));
    return 0;
}
```

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Output:

GCD of 48 and 18 is 6

Explanation:

This program calculates GCD using Euclid's algorithm:

- If $b=0$, GCD is a
- Else GCD is GCD of b and $a\%b$ (remainder of a/b)

4. Write a recursive program to calculate factorial of a number.

Algorithm.

Problem Statement:

Given a non-negative integer n , compute the factorial of n (denoted as $n!$) using a recursive approach.

Factorial Definition:

- $n! = n * (n-1) * (n-2) * \dots * 1$
- $0! = 1$ (by definition)

Recursive Approach:

1. **Base Case:** If n is 0 or 1, return 1 (since $0! = 1$ and $1! = 1$).
2. **Recursive Case:** For $n > 1$, return $n * \text{factorial}(n - 1)$.

Algorithm Steps:

1. Check if $n == 0$ or $n == 1$:
 - If yes, return 1.
2. Otherwise:
 - Return $n * \text{factorial}(n - 1)$.



DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Program:

```
#include <stdio.h>

int fact(int n) {
    if(n == 0) return 1;
    return n * fact(n-1);
}

int main() {
    int n = 5;
    printf("Factorial of %d is %d\n", n, fact(n));
    return 0;
}
```

Output:

Factorial of 5 is 120

Explanation:

- Base case: $0! = 1$
- Recursive case: $n! = n \times (n-1)!$

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

5. Write a program to implement linear searching technique on an array.

Algorithm:

Algorithm

Linear_Search(a, n, val) // 'a' is the given array, 'n' is the size of given array, 'val' is the value to search

Step 1: set **pos** = -1

Step 2: set **i** = 1

Step 3: repeat step 4 while **i** <= n

Step 4: if **a[i]** == val

set **pos** = **i**

print pos

go to step 6

[end of if]

set **ii** = **i** + 1

[end of loop]

Step 5: if **pos** = -1

print "value is not present in the array "

[end of if]

Step 6: exit

Program:

```
#include <stdio.h>
```

```
int linearSearch(int arr[], int n, int key) {
```

```
    for(int i=0; i<n; i++)
```

```
        if(arr[i] == key) return i;
```

```
    return -1;
```

```
}
```

```
int main() {
```

```
    int arr[] = {10, 20, 30, 40, 50};
```

```
    int n = sizeof(arr)/sizeof(arr[0]);
```

```
    int key = 30;
```



DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```
int pos = linearSearch(arr, n, key);  
if(pos != -1) printf("%d found at index %d\n", key, pos);  
else printf("%d not found\n", key);  
  
return 0;  
}
```

Output:

30 found at index 2

Explanation:

Linear search checks each element sequentially until the key is found or end of array is reached.

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

6. Write a program to implement binary searching technique on an array.

Algorithm:

```
Step 1: set beg = lower_bound, end = upper_bound, pos = - 1
Step 2: repeat steps 3 and 4 while beg <=end
Step 3: set mid = (beg + end)/2
Step 4: if a[mid] = val
    set pos = mid
    print pos
    go to step 6
else if a[mid] > val
    set end = mid - 1
else
    set beg = mid + 1
[end of if]
[end of loop]
Step 5: if pos = -1
    print "value is not present in the array"
[end of if]
Step 6: exit
```

Program:

```
#include <stdio.h>

int binarySearch(int arr[], int l, int r, int key) {
    while(l <= r) {
        int m = l + (r-l)/2;
        if(arr[m] == key) return m;
        if(arr[m] < key) l = m+1;
        else r = m-1;
    }
    return -1;
}
```

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```
int main() {  
    int arr[] = {10, 20, 30, 40, 50};  
    int n = sizeof(arr)/sizeof(arr[0]);  
    int key = 40;  
    int pos = binarySearch(arr, 0, n-1, key);  
    if(pos != -1) printf("%d found at index %d\n", key, pos);  
    else printf("%d not found\n", key);  
    return 0;  
}
```

Output:

40 found at index 3

Explanation:

Binary search works on sorted arrays by repeatedly dividing the search interval in half.

7. Write a program to implement a stack using an array. Include the operations Push, Pop, and display the current stack.

Algorithm:

Algorithm

1. Checks if the stack is full.
2. If the stack is full, produces an error and exit.
3. If the stack is not full, increments top to point next empty space.
4. Adds data element to the stack location, where top is pointing.
5. Returns success.

Algorithm

1. Checks if the stack is empty.
2. If the stack is empty, produces an error and exit.
3. If the stack is not empty, accesses the data element at which top is pointing.
4. Decreases the value of top by 1.

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

5. Returns success.

Program:

```
#include <stdio.h>

#define SIZE 5

int stack[SIZE], top = -1;

void push(int val) {
    if(top >= SIZE-1) printf("Stack Overflow!\n");
    else stack[++top] = val;
}

void pop() {
    if(top < 0) printf("Stack Underflow!\n");
    else printf("Popped: %d\n", stack[top--]);
}

void display() {
    if(top < 0) printf("Stack is empty!\n");
    else {
        printf("Stack: ");
        for(int i=0; i<=top; i++) printf("%d ", stack[i]);
        printf("\n");
    }
}

int main() {
    push(10);
    push(20);
    push(30);
    display();
}
```

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```
pop();  
display();  
return 0;  
}
```

Output:

Stack: 10 20 30

Popped: 30

Stack: 10 20

Explanation:

This implements a stack with:

- Push: Adds element at top
- Pop: Removes top element
- Display: Shows stack contents.

PART-B

1. Write a program to implement bubble sort technique.

Algorithm:

```
bubbleSort(array)  
n = length(array)  
repeat  
  swapped = false  
  for i = 1 to n - 1  
    if array[i - 1] > array[i], then  
      swap(array[i - 1], array[i])  
      swapped = true  
    end if  
  end for  
  n = n - 1  
until not swapped  
end bubbleSort
```

Steps of Bubble Sort

Step 1: Compare the first two elements of the input array.

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Step 2: If the first element is greater than the second element, swap them; otherwise, swapping is not required.

Step 3: According to step 2, do the comparison and swapping of the next pair of elements of the input array (the second and the third element), i.e. if the second element is greater than the third element, do the swapping; otherwise, not.

Step 4: Continue this process till the end of the input array is reached. At this point, the largest element will be "bubble up" and get the end of the list.

Step 5: Repeat the steps from 1 to 4, excluding the last element, as it has already achieved its correct position.

Program:

```
#include <stdio.h>

void bubbleSort(int arr[], int n) {
    for(int i=0; i<n-1; i++)
        for(int j=0; j<n-i-1; j++)
            if(arr[j] > arr[j+1]) {
                int temp = arr[j];
                arr[j] = arr[j+1];
                arr[j+1] = temp;
            }
}

int main() {
    int arr[] = {64, 34, 25, 12, 22, 11, 90};
    int n = sizeof(arr)/sizeof(arr[0]);
    bubbleSort(arr, n);
    printf("Sorted array: ");
    for(int i=0; i<n; i++) printf("%d ", arr[i]);

    return 0;
}
```

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

}

Output:

Sorted array: 11 12 22 25 34 64 90

Explanation:

Bubble sort repeatedly swaps adjacent elements if they are in wrong order.

2. Write a program to implement selection sort technique.

Algorithm:

```
// Function to do selection sort on the array 'a' of size 'n'
function selectionSort(a, s)
  // Outer loop: for iterating through the array
  // The last element is not included
  for j = 0 to s - 1
    // Assuming that the current element is the minimum
    // storing its index.
    minIdx = j
    // Inner loop: for finding the minimum element in the unsorted portion
    for k = j + 1 to s - 1
      // If a smaller element is found
      if a[k] < a[minIdx]
        // Updating the index of the element, which is the minimum
        minIdx = k
    // Swapping the minimum element found with the current element
    if minIdx != j
      swap(a[j], a[minIdx])
  // The array is now sorted. return it
  return a
```

Steps of Selection Sort Algorithm

Step 1: Find the smallest element and swap it with the first element of the input array. If the first element is the smallest, then swapping is not required.

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Step 2: Find the smallest element from the remaining elements (the second smallest element) and swap it with the second element. Again, here also swapping is not required if the second smallest element is the second element of the array.

Step 3: Repeat the process for the remaining elements of the input array until all elements achieve the correct position.

Program:

```
#include <stdio.h>

void selectionSort(int arr[], int n) {
    for(int i=0; i<n-1; i++) {
        int min_idx = i;
        for(int j=i+1; j<n; j++)
            if(arr[j] < arr[min_idx]) min_idx = j;

        int temp = arr[min_idx];
        arr[min_idx] = arr[i];
        arr[i] = temp;
    }
}

int main() {
    int arr[] = {64, 25, 12, 22, 11};
    int n = sizeof(arr)/sizeof(arr[0]);
    selectionSort(arr, n);
    printf("Sorted array: ");
    for(int i=0; i<n; i++) printf("%d ", arr[i]);
    return 0;
}
```

Output:

Sorted array: 11 12 22 25 64

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Explanation:

Selection sort finds the minimum element and swaps it with the first unsorted element.

3. Write a program to implement insertion sort technique.

Algorithm:

The simple steps of achieving the insertion sort are listed as follows -

Step 1: If the element is the first element, assume that it is already sorted. Return 1.

Step 2: Pick the next element and store it separately in a **key**.

Step 3: Now, compare the **key** with all elements in the sorted array.

Step 4: If the element in the sorted array is smaller than the current element, then move to the next element. Else, shift greater elements in the array towards the right.

Step 5: Insert the value.

Step 6: Repeat until the array is sorted.

1. INSERTION-SORT(A)
2. for $i \leftarrow 1$ to $\text{length}(A) - 1$ do
3. $\text{key} \leftarrow A[i]$
4. $j \leftarrow i - 1$
5. while $j \geq 0$ and $A[j] > \text{key}$ do
6. $A[j + 1] \leftarrow A[j]$
7. $j \leftarrow j - 1$
8. $A[j + 1] \leftarrow \text{key}$

Program:

```
#include <stdio.h>

void insertionSort(int arr[], int n) {
    for(int i=1; i<n; i++) {
        int key = arr[i];
        int j = i-1;
        while(j >= 0 && arr[j] > key) {
            arr[j+1] = arr[j];
            j--;
        }
    }
}
```

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```
    }  
    arr[j+1] = key;  
  }  
}  
  
int main() {  
    int arr[] = {12, 11, 13, 5, 6};  
    int n = sizeof(arr)/sizeof(arr[0]);  
    insertionSort(arr, n);  
    printf("Sorted array: ");  
    for(int i=0; i<n; i++) printf("%d ", arr[i]);  
    return 0;  
}
```

Output:

Sorted array: 5 6 11 12 13

Explanation:

Insertion sort builds the final array one item at a time by inserting each element in its correct position.

4. Write a program to convert an infix expression to a postfix expression using stack.

Algorithm:

Step 1. Scan the expression from left to right

Step 2. If it is an operand push it to stack

Step 3. If it is an operator pull operand from stack and perform operation

Step 4. Store the output of step 3, back to stack

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Step 5. Scan the expression until all operands are consumed

Step 6. Pop the stack and perform operation.

Program:

```
#include <stdio.h>
#include <ctype.h>
#include <string.h>
#define SIZE 50

char stack[SIZE];
int top = -1;

void push(char elem) { stack[++top] = elem; }
char pop() { return stack[top--]; }
int precedence(char op) {
    Switch(op) {
        Case '+': case '-': return 1;
        Case '*': case '/': return 2;
        Case '^': return 3;
    }
    Return 0;
}

Void infixToPostfix(char infix[], char postfix[]) {
    Int i = 0, j = 0;
    Char ch, x;
    Push("(");
    Strcat(infix, ")");
```



DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```

While((ch = infix[i++]) != '\0') {
    If(ch == '(') push(ch);
    Else if(isalnum(ch)) postfix[j++] = ch;
    Else if(ch == ')') {
        While((x = pop()) != '(') postfix[j++] = x;
    }
    Else {
        While(precedence(stack[top]) >= precedence(ch))
            Postfix[j++] = pop();
        Push(ch);
    }
}
Postfix[j] = '\0';
}

int main() {
    Char infix[SIZE] = "a+b*(c^d-e)^(f+g*h)-I";
    Char postfix[SIZE];

    infixToPostfix(infix, postfix);
    printf("Infix: %s\nPostfix: %s\n", infix, postfix);
    return 0;
}

```

Output:

Infix: a+b*(c^d-e)^(f+g*h)-i

Postfix: abcd^e-fgh*+^*+i-

Explanation:

This converts infix expressions to postfix notation using stack to handle operator precedence.

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

5. Write a program to implement an ordinary queue using an array. Include the operations Enqueue, Dequeue, and Display.

Algorithm:

Queue Operations:

1. **Enqueue(item):** Add an item to the rear of the queue.
2. **Dequeue():** Remove and return the item from the front of the queue.
3. **Display():** Print all elements in the queue from front to rear.

Data Structures Used:

- An array `queue[]` to store elements.
- Two pointers:
 - `front` (index of the front element).
 - `rear` (index of the rear element).
- A constant `MAX_SIZE` to define the maximum capacity of the queue.

Algorithm Steps:

1. Initialize the Queue:

- Set `front = -1` and `rear = -1` (empty queue).

2. Enqueue(item):

1. Check if the queue is full (`rear == MAX_SIZE - 1`):
 - If full, print "Queue Overflow" and return.
2. If the queue is empty (`front == -1`), set `front = 0`.
3. Increment `rear` by 1.
4. Insert `item` at `queue[rear]`.

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

3. Dequeue():

1. Check if the queue is empty (`front == -1` or `front > rear`):
 - If empty, print "Queue Underflow" and return `None`.
2. Store `queue[front]` in a temporary variable.
3. Increment `front` by 1.
4. If `front > rear` (queue becomes empty), reset `front = rear = -1`.
5. Return the stored item.

4. Display():

1. Check if the queue is empty (`front == -1` or `front > rear`):
 - If empty, print "Queue is empty".
2. Else, loop from `front` to `rear` and print each element.

Program:

```
#include <stdio.h>
```

```
#define SIZE 5
```

```
int queue[SIZE], front = -1, rear = -1;
```

```
void enqueue(int val) {
```

```
    if(rear == SIZE-1) printf("Queue is full!\n");
```

```
    else {
```

```
        if(front == -1) front = 0;
```

```
        Queue[++rear] = val;
```

```
    }
```

```
}
```

```
void dequeue() {
```

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```
if(front == -1 || front > rear) printf("Queue is empty!\n");  
else printf("Dequeued: %d\n", queue[front++]);  
}
```

```
void display() {  
    if(front == -1) printf("Queue is empty!\n");  
    else {  
        Printf("Queue: ");  
        For(int i=front; i<=rear; i++) printf("%d ", queue[i]);  
        Printf("\n");  
    }  
}
```

```
int main() {  
    Enqueue(10);  
    Enqueue(20);  
    Enqueue(30);  
    Display();  
    Dequeue();  
    Display();  
  
    Return 0;  
}
```

Output:

Queue: 10 20 30

Dequeued: 10

Queue: 20 30

Explanation:

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

This implements a queue with:

- Enqueue: Adds element at rear
- Dequeue: Removes element from front
- Display: Shows queue contents

6. Write a program to create a singly linked list and perform the following operations: Insertion at the beginning, Deletion from the beginning, and Traversal of the linked list.

Algorithm:

1. Node Structure

- Each node has:
 - data: Stores the value.
 - next: Pointer to the next node.

2. Linked List Operations

- **Insertion at Beginning (insert_at_beginning)**
 1. Create a new node with the given data.
 2. Point the new node's next to the current head.
 3. Update head to the new node.
- **Deletion from Beginning (delete_from_beginning)**
 1. Check if the list is empty (head == None). If yes, return.
 2. Store the head node in a temporary variable.
 3. Move head to head.next.
 4. Free the temporary node (optional in Python due to garbage collection).
- **Traversal (display)**
 1. Start from head.
 2. Traverse each node while printing its data.
 3. Stop when current.next == None.



DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Program:

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
Struct Node {  
    Int data;  
    Struct Node *next;  
};
```

```
Struct Node *head = NULL;
```

```
Void insertAtBeginning(int val) {  
    Struct Node *newNode = (struct Node*)malloc(sizeof(struct Node));  
    newNode->data = val;  
    newNode->next = head;  
    head = newNode;  
}
```

```
Void deleteFromBeginning() {  
    If(head == NULL) printf("List is empty!\n");  
    Else {  
        Struct Node *temp = head;  
        Head = head->next;  
        Printf("Deleted: %d\n", temp->data);  
        Free(temp);  
    }  
}
```



DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```
Void display() {  
    Struct Node *current = head;  
    Printf("List: ");  
    While(current != NULL) {  
        Printf("%d ", current->data);  
        Current = current->next;  
    }  
    Printf("\n");  
}
```

```
Int main() {  
    insertAtBeginning(30);  
    insertAtBeginning(20);  
    insertAtBeginning(10);  
    display();  
  
    deleteFromBeginning();  
    display();  
  
    return 0;  
}
```

Output:

List: 10 20 30

Deleted: 10

List: 20 30

Explanation:

This implements a singly linked list with:

- Insertion at beginning

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

- Deletion from beginning
- Traversal to display list

7. Write a program to represent a binary tree and implement Pre-order ,In-order and Post-order traversal methods.

Algorithm:

1. Binary Tree Node Structure

- Each node contains:
 - data (integer value)
 - left (pointer to the left child)
 - right (pointer to the right child)

2. Tree Traversal Methods

1. Pre-order Traversal (Root → Left → Right)

- Visit the root.
- Traverse the left subtree.
- Traverse the right subtree.

2. In-order Traversal (Left → Root → Right)

- Traverse the left subtree.
- Visit the root.
- Traverse the right subtree.

3. Post-order Traversal (Left → Right → Root)

- Traverse the left subtree.
- Traverse the right subtree.
- Visit the root.

DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

Program:

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
Struct Node {  
    Int data;  
    Struct Node *left, *right;  
};
```

```
Struct Node* createNode(int val) {  
    Struct Node* newNode = (struct Node*)malloc(sizeof(struct Node));  
    newNode->data = val;  
    newNode->left = newNode->right = NULL;  
    return newNode;  
}
```

```
Void preorder(struct Node* root) {  
    If(root == NULL) return;  
    Printf("%d ", root->data);  
    Preorder(root->left);  
    Preorder(root->right);  
}
```

```
Void inorder(struct Node* root) {  
    If(root == NULL) return;  
    Inorder(root->left);  
    Printf("%d ", root->data);  
    Inorder(root->right);  
}
```



DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

```
}  
  
Void postorder(struct Node* root) {  
    If(root == NULL) return;  
    Postorder(root->left);  
    Postorder(root->right);  
    Printf("%d ", root->data);  
}  
  
Int main() {  
    Struct Node* root = createNode(1);  
    Root->left = createNode(2);  
    Root->right = createNode(3);  
    Root->left->left = createNode(4);  
    Root->left->right = createNode(5);  
  
    Printf("Preorder: "); preorder(root); printf("\n");  
    Printf("Inorder: "); inorder(root); printf("\n");  
    Printf("Postorder: "); postorder(root); printf("\n");  
  
    Return 0;  
}
```

Output:

Preorder: 1 2 4 5 3

Inorder: 4 2 5 1 3

Postorder: 4 5 2 3 1

Explanation:

This demonstrates three tree traversal methods:

- Preorder: Root, Left, Right



DEPARTMENT OF BACHELOR OF COMPUTER APPLICATIONS

- Inorder: Left, Root, Right
- Postorder: Left, Right, Root